On The Delays in Spiking Neural P Systems *

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ABSTRACT

In this work we extend and improve the results obtained in a previous work on simulating Spiking Neural P systems (SNP systems in short) with delays using SNP systems without delays. We simulate the former with the latter over sequential, iteration, join, and split routing. Our results provide constructions so that both systems halt at exactly the same time, start with only one spike, and produce the same number of spikes to the environment after halting.

Keywords: Membrane Computing, Spiking Neural P systems, delays, simulation, routing

1. INTRODUCTION

SNP systems, first presented in [7] and with some recent results in [12], [13] and [10] (among others), are computing devices inspired by how biological neurons represent information: using electro-chemical signals called spikes. Since spikes are indistinct, information is taken not from the spikes themselves, but from their multiplicity or time of arrival. One motivation for SNP systems (as is the case in the area of Membrane Computing [11] in general) is to abstract ideas from biology for computational use. For SNP systems in particular, the neuron from our brains is the motivation. It can be argued that the human brain is one (if not currently) the most complicated and powerful "supercomputer" known to us at the moment. The brain performs complex computations from billions of interconnected neurons while consuming only around 10 to 20 Watts of energy [8], and it is small enogh to fit in our skulls. It is therefore desirable to work with as little quantity of "energy" as possible, and we can think of the spike in SNP systems as being such quantity.

SNP systems are Turing complete devices [7, 5] and have been used as (among others) transducers [6], generating vectors of numbers [1], as well solving hard problems [9]. Spiking rules (rules that produce spikes) are usually of two types: with delays and without delays. If an SNP system has at least one rule with a delay, we refer to it as an SNP system with delay labeled as Π , otherwise it is known as an SNP system without delay labeled as $\overline{\Pi}$. In [6] it was shown that SNP systems without delay are Turing complete i.e. delays are not required for computational completeness. However (and as with other models of computation that still incorporate time in their variants) SNP systems with delays still provide improvements over those without delays e.g. a Π is more compact i.e. has less neurons, compared to $\overline{\Pi}$ that simulates Π (related to their descriptional complexity).

In this work we extend the work presented in [3], with the goal of simulating a Π that performs sequential, iteration, join, and split routing with a $\overline{\Pi}$ that performs the same routings. By routing we mean the transfer or movement of spikes from one neuron to another. By simulation in [3] it is meant that the following two requirements are satisfied:

 ${f R_1}$: Halting time of Π coincides with the halting time of $\overline{\Pi}$, or is offset either by a fixed timestep or by a function of the delays in Π ,

 $\mathbf{R_2}$: number of spikes in the final configuration of Π is the same number in $\overline{\Pi}$, or is offset by a function of the delays in Π .

In [3], the construction of $\overline{\Pi}$ from Π is such that the initial spikes of $\overline{\Pi}$ is a function of the delay (or delays) in Π . In particular, the initial spikes of $\overline{\Pi}$ are multiples of the delays in Π . Aside from the increased initial spike number in $\overline{\Pi}$, the exponents in the regular expressions and the number of consumed spikes of certain spiking rules in $\overline{\Pi}$ are also multiples of all the delays in a given routing of Π .

We improve the work done in [3] by providing alternative constructions in this work. Our specific contributions are as follows:

• we construct a $\overline{\Pi}$ that simulates a Π that performs sequential, iteration, join, and split routing,

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- both $\overline{\Pi}$ and Π start with only one spike each in the initial neuron.
- halting time of $\overline{\Pi}$ and Π coincide i.e. there are no offsets,
- number of spikes sent to the environment after halting are equal for Π and Π̄.
- our construction allows split routing even if the delays of the output neurons are not equal.

The trade-off is that for every delay d in Π , we add d neurons in $\overline{\Pi}$. If the $initial\ neuron(s)$ (i.e. the first set of neuron(s) to spike in the computation) of Π has a delay, following our construction means we simply modify Π and $\overline{\Pi}$ such that their new halting time involves one additional time step. The succeeding sections are as follows: Section 2 provides preliminaries and assumptions for our work. Section 3 presents our main results. We end with our final remarks and directions for future work in Section 4.

2. PRELIMINARIES

It is assumed that the readers are familiar with the basics of Membrane Computing (a good introduction is [11] with recent results and information in the P systems webpage at http://ppage.psystems.eu/ and a recent handbook in [13]) and formal language theory. We only briefly mention notions and notations which will be useful throughout the paper.

Let V be an alphabet, V^* is the free monoid over V with respect to concatenation and the identity element λ (the empty string). The set of all non-empty strings over V is denoted as V^+ so $V^+ = V^* - \{\lambda\}$. We call V a singleton if $V = \{a\}$ and simply write a^* and a^+ instead of $\{a^*\}$ and $\{a^+\}$. The length of a string $w \in V^*$ is denoted by |w|. If a is a symbol in V, $a^0 = \lambda$. A language $L \subseteq V^*$ is regular if there is a regular expression E over V such that L(E) = L. A regular expression over an alphabet V is constructed starting from λ and the symbols of V using the operations union, concatenation, and +, using parentheses when necessary to specify the order of operations. Specifically, (i) λ and each $a \in V$ are regular expressions, (ii) if E_1 and E_2 are regular expressions over V then $(E_1 \cup E_2)$, E_1E_2 , and E_1^+ are regular expressions over V, and (iii) nothing else is a regular expression over V. With each expression E we associate a language L(E) defined in the following way: (i) $L(\lambda) = \{\lambda\}$ and $L(a) = \{a\}$ for all $a \in V$, (ii) $L(E_1 \cup E_2) = L(E_1) \cup L(E_2), L(E_1 E_2) = L(E_1)L(E_2), \text{ and }$ $L(E_1^+) = L(E_1)^+$, for all regular expressions E_1 , E_2 over V. Unnecessary parentheses are omitted when writing regular expressions, and $E^+ \cup \{\lambda\}$ is written as E^* . Next we have the definition for an SNP system.

DEFINITION 1 (SNP SYSTEM). An SNP system of a finite degree $m \ge 1$ is a construct of the form

$$\Pi = (O, \sigma_1, \ldots, \sigma_m, syn, out),$$

where:

1. $O = \{a\}$ is the singleton alphabet (a is called spike).

- 2. $\sigma_1, \ldots, \sigma_m$ are neurons of the form $\sigma_i = (n_i, R_i), 1 \le i \le m$, where:
 - (a) $n_i \geq 0$ is an integer representing the number of spikes in σ_i
 - (b) R_i is a finite set of rules of the general form

$$E/a^c \rightarrow a^b; d$$

where E is a regular expression over O, $c \ge 1$, if b > 0 then $d \ge 0$ and $c \ge b$, else if b = 0 then d = 0.

- 3. $syn \subseteq \{1, 2, ..., m\} \times \{1, 2, ..., m\}, (i, i) \notin syn for 1 \le i \le m, are synapses between neurons.$
- 4. $out \in \{1, 2, ..., m\}$ is the index of the output neuron.

A spiking rule is where $b \geq 1$. A forgetting rule is a rule where b = 0 is written as $E/a^c \to \lambda$. If $L(E) = \{a^c\}$ then spiking and forgetting rules are simply written as $a^c \to a^b$ and $a^c \to \lambda$, respectively. Applications of rules are as follows: if neuron σ_i contains k spikes, $a^k \in L(E)$ and $k \geq c$, then the rule $E/a^c \to a^b \in R_i$ is enabled and the rule can be fired or applied. If $b \geq 1$, the application of this rule removes c spikes from σ_i , so that only k-c spikes remain in σ_i . The neuron sends b number of spikes to every σ_j such that $(i,j) \in syn$. The output neuron has a synapse not directed to any other neuron, only to the environment. The neuron σ_1 is referred to as the initial neuron.

If a spiking rule (forgetting rules cannot have delays) has d=0, the b number of spikes are sent immediately i.e. in the same time step as the application of the rule. If $d\geq 1$ and the spiking rule was applied at time t, then the spikes are sent at time t+d. From time t to t+d-1 the neuron is said to be closed (inspired by the $refractory\ period$ of the neuron in biology) and cannot receive spikes. Any spikes sent to the neuron when the neuron is closed are lost or removed from the system. At time t+d the neuron becomes open and can then receive spikes again. The neuron can then apply another rule at time t+d+1. If b=0 then no spikes are produced. SNP systems assume a global clock, so the application of rules and the sending of spikes by neurons are all synchronized.

A configuration of the system at time k is denoted as $C_k = \langle n_1/t_1, \ldots, n_m/t_m, n_e \rangle$, where each element of the vector (except for n_e , denoting the spikes in the environment) is the configuration of a neuron σ_i , with n_i spikes and is open after $t_i \geq 0$ steps. An initial configuration C_0 is therefore $\langle n_1/0, \ldots, n_m/0, 0 \rangle$ since no rules whether with or without delay, have yet been applied and the environment is initially empty. A computation is a sequence of transitions from an initial configuration. A computation may halt (no more rules can be applied for a given configuration) or not. If an SNP system does halt, all neurons should be open. Computation result in this work is obtained by checking the number of spikes in the environment once the system halts.

As an example, let us have an SNP system shown in Figure 1 formally defined as follows: $\Pi_0 = (O, \sigma_1, \sigma_2, \sigma_3, syn, out)$ where $\sigma_1 = (1, a^+/a \rightarrow a), \ \sigma_2 = (0, a^+/a \rightarrow a; 2), \ \sigma_3 =$

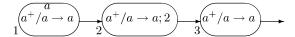


Figure 1: SNP system with delay Π_0 .

 $(0,a^+/a \to a)$, $syn = \{(1,2), (2,3)\}$, the source neuron is σ_1 , and the output neuron is σ_3 . Only neuron σ_1 has one spike at the beginning of the computation and only σ_2 has a rule with a delay d=2. We have $C_0 = \langle 1/0, 0/0, 0/0, 0\rangle$. At the next step, σ_1 can use its rule (it has at least one spike) and consumes one spike and sends one spike immediately to σ_2 so we have $C_1 = \langle 0/0, 1/0, 0/0, 0\rangle$. At step 2, σ_2 consumes its spike and closes for 2 time steps, so $C_2 = \langle 0/0, 0/2, 0/0, 0\rangle$. At step 3 we have $C_3 = \langle 0/0, 0/1, 0/0, 0\rangle$. At time step 4, σ_2 opens and sends one spike to σ_3 , so $C_3 = \langle 0/0, 0/0, 1/0, 0\rangle$. Finally, at time step 5 the output neuron sends one spike to the environment, Π_0 halts and we have $C_5 = \langle 0/0, 0/0, 0/0, 1/0, 1\rangle$.

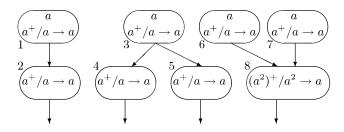


Figure 2: Routing constructs (from left to right): sequential, split, and join.

SNP systems where each neuron has exactly one rule are called simple, while the systems that have the same set of rules are called homogeneous [14]. In this work, if SNP systems only have rules of the restricted form $(a^k)^+/a^k \rightarrow a$ where k is a non-negative integer, we refer to them as semihomogeneous. We only consider SNP systems Π and $\overline{\Pi}$ that are simple and semi-homogeneous, where their initial configurations have one spike in the initial neuron only, and no spike in every other neuron (as in Figure 1) in this work. We make no restrictions on the values of the delays in a (rule of a) neuron. The objective is to route or move the single spike in the initial neuron through the system, towards the output neuron, and eventually to the environment. Spikes are routed via paths, where a path consists of at least two neurons σ_i, σ_j such that $(i, j) \in syn$. Using paths, we can have four basic routing constructs (referring to Figure 2):

- 1. sequential where, given at least two neurons σ_1 , σ_2 such that σ_2 spikes only after σ_1 spikes and there is a path from σ_1 to σ_2 ,
- 2. iteration, where at least two neurons spike multiple (possibly an infinite) number of times and a loop is formed e.g. adding a synapse (2,1) which creates a loop between σ_1 and σ_2 ,
- 3. *split*, where a spike from σ_3 is sent to at least two output neurons σ_4 and σ_5 and $(3,4),(3,5) \in syn$,

4. *join*, where spikes from at least two input neurons σ_6, σ_7 are sent to a neuron σ_8 , where $(6,8), (7,8) \in syn$, so that σ_8 produces a spike only after accumulating spikes from σ_7 and σ_8 .

Notice that iteration routing can be formed by combining the three other constructs. Also notice that if there exists a sequential path from σ_i (with delay d_1) to σ_j (with delay d_2) so that $d_1 < d_2$ and the number of spikes of the initial neuron σ_1 in C_0 is $n_1 > 1$, it is possible for some spikes to be lost. The reason is that it is possible for σ_j to still be closed when spikes from σ_i arrive. We avoid lost spikes by considering SNP systems where the initial neuron has only one spike. We say in this work that a $\overline{\Pi}$ simulates a Π if two requirements are satisfied:

 \mathbf{R}'_1 : halting time of Π spikes is the same halting time of $\overline{\Pi}$,

 $\mathbf{R'_2}$: number of spikes in the environment of Π when Π halts is equal to the number of spikes in the environment of $\overline{\Pi}$ when $\overline{\Pi}$ halts.

3. MAIN RESULTS

We begin presenting our results with a fundamental idea on sequential routing.

Lemma 1 (Sequential routing). Given an SNP system with delay Π performing sequential routing, there exists an SNP system without delay $\overline{\Pi}$ performing sequential routing that simulates Π .

PROOF. We refer to Figure 3 for illustrations. Let

$$\Pi = (O, \sigma_{11}, \sigma_{12}, \{(11, 12)\}, 12)$$
 with $\sigma_1 = (1, a^+/a \rightarrow a)$ and $\sigma_2 = (0, a^+/a \rightarrow a; d)$

we then let

 $\overline{\Pi} = (O, \sigma_{21}, \sigma_{22-i}, \sigma_{23}, syn, 23)$ where $1 \le i \le d$, $syn = \{(21, 22-1), \dots, (21, 22-(d-1)), (22-1, 22-d), \dots, (22-(d-1), 22-d), (22-d, 23)\}, \sigma_{21} = (1, a^+/a \to a), \sigma_{22-i} = (0, a^+/a \to a), \sigma_{23} = (0, (a^{d-1})^+/a^{d-1} \to a).$

The additional d neurons, immediately after initial neuron σ_{21} in $\overline{\Pi}$, are used to multiply the single spike from σ_{21} . The additional neurons then send one spike each to σ_{22-d} . Neuron σ_{22-d} accumulates d-1 spikes, and consumes these, one spike at a time and sending one spike every time to σ_{23} . This consumption of one spike every time step creates a delay of d-1 time steps. Due to the regular expression of the rule in σ_{23} , the neuron will have to accumulate d-1 spikes before the rule is used. Once σ_{23} accumulates d-1 spikes, it immediately sends one spike to the environment. This spiking and halting occurs at time t+d+1 for both Π and $\overline{\Pi}$ (satisfying \mathbf{R}'_1 and \mathbf{R}'_2) if we let t be the time when σ_{11} and σ_{21} spike.

We can repeatedly apply the previous construction if there exist more than one neuron with a (rule having a) delay in a sequential path as seen in Figure 4. It can be easily shown that if there exists σ_i without delay in a sequential

path between σ_{11} and σ_{12} , the time to halt for both Π_1 and $\overline{\Pi}_1$ still coincide. In particular, every additional σ_i having a rule without a delay adds one time step to the halting time of both Π and $\overline{\Pi}$. Both \mathbf{R}'_1 and \mathbf{R}'_2 are still satisfied. \square

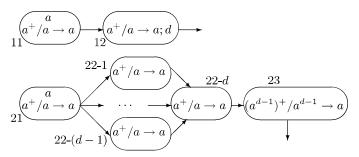


Figure 3: Sequential routing: Π_1 (top) with delay d, and $\overline{\Pi}_1$ (bottom) simulating Π_1 .

A sample computation of Π_1 and $\overline{\Pi}_1$ is shown in Table 1. For sample computations of Π_2 and $\overline{\Pi}_2$ we refer to Table 2. From Lemma 1 we have the following observation.

OBSERVATION 1. If Π has more than one neuron with a delay in a rule, the total additional neurons in $\overline{\Pi}$ is $\sum_{i=1}^{m} d_i$ where d_i is the delay of the rule in σ_i .

Steps	Π_1	$\overline{\Pi}_1$
t_0	$\langle 1/0, 0/0, 0 \rangle$	$\langle 1,0,0,0,0,0 \rangle$
t_1	$\langle 0/0, 1/0, 0 \rangle$	$\langle 0, 1, 1, 0, 0, 0 \rangle$
t_2	$\langle 0/0,0/3,0\rangle$	$\langle 0, 0, 0, 2, 0, 0 \rangle$
t_3	$\langle 0/0,0/2,0\rangle$	$\langle 0,0,0,0,1,0 \rangle$
t_4	$\langle 0/0,0/1,0\rangle$	$\langle 0,0,0,0,2,0 \rangle$
t_5	$\langle 0/0,0/0,1\rangle$	$\langle 0,0,0,0,0,1 \rangle$

Table 1: Sample computations of Π_1 and $\overline{\Pi}_1$, d=3.

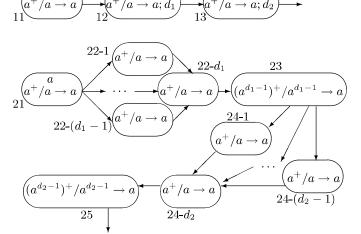


Figure 4: Sequential routing with multiple delays: Π_2 (top) with delays d_1 and d_2 , and $\overline{\Pi}_2$ (bottom) simulating Π_2 .

Steps	Π_2	$\overline{\Pi}_2$
t_0	$\langle 1/0, 0/0, 0/0, 0 \rangle$	$\langle 1, 0, 0, 0, 0, 0, 0, 0, 0, 0 \rangle$
t_1	$\langle 0/0, 1/0, 0/0, 0 \rangle$	$\langle 0, 1, 0, 0, 0, 0, 0, 0, 0, 0 \rangle$
t_2	$\langle 0/0, 0/2, 0/0, 0 \rangle$	$\langle 0, 0, 1, 0, 0, 0, 0, 0, 0, 0 \rangle$
t_3	$\langle 0/0, 0/1, 0/0, 0 \rangle$	$\langle 0, 0, 0, 1, 0, 0, 0, 0, 0, 0 \rangle$
t_4	$\langle 0/0, 0/0, 1/0, 0 \rangle$	$\langle 0, 0, 0, 0, 1, 1, 0, 0, 0 \rangle$
t_5	(0/0, 0/0, 0/3, 0)	$\langle 0, 0, 0, 0, 0, 0, 2, 0, 0 \rangle$
t_6	$\langle 0/0, 0/0, 0/2, 0 \rangle$	$\langle 0, 0, 0, 0, 0, 0, 1, 1, 0 \rangle$
t_7	$\langle 0/0, 0/0, 0/1, 0 \rangle$	$\langle 0, 0, 0, 0, 0, 0, 0, 2, 0 \rangle$
t_8	$\langle 0/0, 0/0, 0/0, 1 \rangle$	$\langle 0, 0, 0, 0, 0, 0, 0, 0, 1 \rangle$

Table 2: Sample computations of Π_2 and $\overline{\Pi}_2$, $d_1=2$, $d_2=3$.

Lemma 2 (Iteration routing). Given an SNP system with delay Π performing iteration routing, there exists an SNP system without delay $\overline{\Pi}$ performing iteration routing that simulates Π .

PROOF. We refer to Figure 5 for illustrations. Let

$$\Pi = (O, \sigma_{11}, \sigma_{12}, \{(11, 12), (12, 11)\}, 12)$$
 where $\sigma_{11} = (1, a^+/a \rightarrow a; d, \sigma_2 = (0, a^+/a \rightarrow a)$

we then let

The construction of $\overline{\Pi}$ uses the construction idea in Lemma 1 i.e. the neuron with a delay d in Π is replaced with d additional neurons in $\overline{\Pi}$. In Figure 5 an infinite loop is created: a spike starts at σ_{11} and it uses its rule at time t so that the spike is sent to σ_{12} at time t+d, then σ_{12} immediately sends a spike back to σ_{11} (and the environment) at time t+d+1, and so on and so forth. Similarly, σ_{21} sends a spike to neurons σ_{22-1} to $\sigma_{22-(d-1)}$ at time t. At time t+1, σ_{22-d} accumulates d-1 spikes from the d-1 neurons from the previous time step. The spikes in σ_{22-d} are consumed and then sent one at a time to σ_{23} . At time t+d, σ_{23} accumulates d-1 spikes so that it sends one spike back to σ_{21} and at the environment at time t+d+1, coinciding with the time of spiking of σ_{12} . Thus, \mathbf{R}'_1 and \mathbf{R}'_2 are satisfied. \square

Lemma 2 for iteration routing makes use of the construction used in Lemma 1 for sequential routing. This construction will again be used for the join and split routings as follows. From Lemma 2 we have the following observation.

Observation 2. If the initial neuron of Π has a delay and its halting time is t+d, we add a new initial neuron $\sigma_{1'}$ in Π with $(1',1) \in syn$ so that Π halts at time t+d+1. We then add a new initial neuron similarly to $\overline{\Pi}$ and modify its syn (following Lemma 1 construction) so that $\overline{\Pi}$ halts at t+d+1 instead, simulating Π .

Although the premise of Observation 2 is different from our assumption in Section 2 that the initial neuron has no delay, the observation provides a solution on how to approach

such a premise. For example, if σ_{11} has a delay instead of σ_{23} in Π_3 , we add a new initial neuron to $\sigma_{11'}$ and a new synapse $(11',11) \in syn$ in Π . For $\overline{\Pi}_3$, we modify it as follows: add a new initial neuron $\sigma_{21'}$ and σ_{21} is replaced with d neurons (instead of σ_{22}). The synapse set of $\overline{\Pi}$ is changed to $\{(21',21-1),\ldots,(21',21-(d-1)),(21-1,21-d),\ldots,(21-(d-1),21-d),(21-d,22),(22,21-1),\ldots,(22,21-(d-1))\}$, we remove σ_{23} and have σ_{22} as the output neuron instead. Both Π_3 and $\overline{\Pi}_3$ halt at the same time at t+d+2.

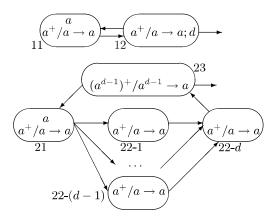


Figure 5: Iteration routing: Π_3 (top) has a delay d simulated by $\overline{\Pi}_3$ (bottom).

LEMMA 3 (JOIN ROUTING). Given an SNP system with delay Π performing join routing, there exists an SNP system without delay $\overline{\Pi}$ performing join routing that simulates Π .

PROOF. We refer to Figure 6 for illustrations. Let

$$\Pi = (O, \sigma_{11}, \sigma_{12}, \sigma_{13}, \{(11, 13), (12, 13)\}, 13)$$
 where $\sigma_{11} = \sigma_{12} = (1, a^+/a \to a), \sigma_{13} = (0, (a^2)^+/a^2 \to a; d)$

we then let

 $\overline{\Pi} = (O, \sigma_{21}, \sigma_{22}, \sigma_{23\text{-}i}, \sigma_{24}, syn, 24) \text{ where } 1 \leq i \leq d, \sigma_{11} = \sigma_{22} = (1, a^+/a \to a), \sigma_{23\text{-}1} = \dots = \sigma_{23\text{-}(d-1)} = (0, (a^2)^+/a^2 \to a^2), \ \sigma_{23\text{-}d} = (0, (a^2)^+/a^2 \to a), \ syn = \{(21, 23\text{-}1), \dots, (21, 23\text{-}(d-1)), (22, 23\text{-}1), \dots, (22, 23\text{-}(d-1)), (23\text{-}1, 23\text{-}d), \dots, (23\text{-}(d-1), 23\text{-}d), (23\text{-}d, 24)\}.$

For Π and $\overline{\Pi}$ we have as initial neurons σ_{11}, σ_{12} and σ_{21}, σ_{22} respectively. Using the construction in Lemma 1, $\overline{\Pi}$ has d additional neurons corresponding to σ_{13} in Π . Let time t be the time when the initial neurons spike. At time t, neurons $\sigma_{23\text{--}1}$ to $\sigma_{23\text{--}(d-1)}$ have two spikes each, so that in total, $\overline{\Pi}$ at this time has 2(d-1) spikes. At the next time step t+1, the d-1 neurons send two spikes each to σ_{23-d} so that σ_{23-d} accumulates 2(d-1) spikes. Since σ_{23-d} consumes two spikes every time and it has 2(d-1) spikes, σ_{23-d} will take d-1time steps to consume all of its 2(d-1) spikes. Every time σ_{23-d} spikes, it sends only one spike to σ_{24} . At time t+d, σ_{24} has accumulated d-1 spikes from σ_{23-d} so that σ_{24} sends one spike to the environment and halts at t + d + 1. This time step coincides with the halting time of σ_{13} in Π , sending one spike to the environment. We therefore satisfy $\mathbf{R_1'}$ and $\mathbf{R_2'}$.

A sample computation of Π_4 and $\overline{\Pi}_4$ is shown in Table 3.

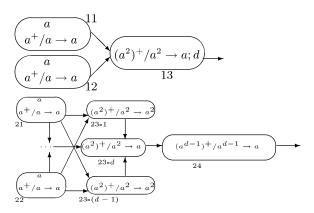


Figure 6: Join routing : Π_4 (top) has delay d simulated by $\overline{\Pi}_4$ (bottom).

Steps	Π_4	$\overline{\Pi}_4$
t_0	$\langle 1/0, 1/0, 0/0, 0 \rangle$	$\langle 1, 1, 0, 0, 0, 0, 0, 0 \rangle$
t_1	$\langle 0/0,0/0,2/0,0\rangle$	$\langle 0,0,2,2,0,0,0 \rangle$
t_2	$\langle 0/0,0/0,0/3,0\rangle$	$\langle 0,0,0,0,4,0,0 \rangle$
t_3	$\langle 0/0,0/0,0/2,0\rangle$	$\langle 0,0,0,0,2,1,0 \rangle$
t_4	$\langle 0/0,0/0,0/1,0\rangle$	$\langle 0,0,0,0,0,2,0 \rangle$
t_5	$\langle 0/0,0/0,0/0,1\rangle$	(0,0,0,0,0,0,1)

Table 3: Sample computations of Π_4 and $\overline{\Pi}_4$, d=3.

Lemma 4 (Split routing). Given an SNP system with delay Π performing split routing, there exists an SNP system without delay $\overline{\Pi}$ performing split routing that simulates Π .

PROOF. We refer back to the split routing in Figure 2. Notice that a split routing can be thought of as two separate paths, either from σ_3 to σ_4 or σ_3 to σ_5 . We let t be the time that σ_3 spikes and modify the split routing in Figure 2 as follows and let it be

$$\Pi = (O, \sigma_3, \sigma_4, \sigma_5, \sigma_o, syn, o)$$
 where $\sigma_3 = (1, a^+/a \rightarrow a), \sigma_5 = (0, a^+/a \rightarrow a), \sigma_4 = (0, a^+/a \rightarrow a; d), syn = \{(3, 4), (3, 5), (4, o), (5, o)\}.$

We arbitrarily chose σ_4 to have a delay instead of σ_5 in this case. Next we let

$$\overline{\Pi} = (O, \sigma_{3'}, \sigma_{4'-i}, \sigma_{5'}, \sigma_o, syn, o) \text{ where } 1 \leq i \leq d, \sigma_{3'} = (1, a^+/a \to a), \sigma_{4'-1} = \dots = \sigma_{4'-(d-1)} = \sigma_{5'} = (0, a^+/a \to a), \sigma_{4'-d} = (0, (a^{d-1})^+/a^{d-1} \to a), syn = \{(3', 4'-1), \dots, (3', 4'-(d-1)), (3', 5'), (4'-1, 4'-d), \dots, (4'-(d-1), 4'-d), (4'-d, o), (5, o)\}.$$

Since σ_4 has delay d, we simply follow Lemma 1 and add d neurons in $\overline{\Pi}$ corresponding to σ_4 . Let t be the time when σ_3 and $\sigma_{3'}$ spike. The time that σ_o spikes the second time (since the spike from σ_5 makes σ_o spike the first time, followed by the delayed spike from σ_4) i.e. the halting time, is t+d+1 and the environment receives two spikes in total. This time coincides with the halting time of $\overline{\Pi}$, which also sends two spikes to its environment. \mathbf{R}'_1 and \mathbf{R}'_2 are both satisfied for this case.

In the case where both σ_4 and σ_5 have delays d_4 and d_5

respectively, then following Lemma 1, $\overline{\Pi}$ has $d_4 + d_5$ additional neurons. For both systems, halting time is $t + d_{max} + 1$ where $d_{max} = max(d_4, d_5)$, and two spikes are sent to the environment, satisfying \mathbf{R}'_1 and \mathbf{R}'_2 . \square

We can now have the following theorem.

Theorem 1. Given an SNP system Π with delays containing one or more of the following routings: sequential, iteration, join, split, there exists an SNP system $\overline{\Pi}$ that simulates Π .

PROOF. Proof follows from Lemma 1, 2, 3, 4. \Box

Notice that Observation 1 and 2 hold for all four routing constructs.

4. FINAL REMARKS

We have presented an alternative construction of a $\overline{\Pi}$ that simulates a given Π , improving the previous work so that we use only one initial spike for both systems. The halting time of both systems also exactly coincide with one another. The trade off is that there is an "explosion" of neurons in $\overline{\Pi}$ for every delay in Π i.e. we add d_i neurons in $\overline{\Pi}$ for every σ_i that has a delay.

For our further work, we will consider nondeterministic SNP systems i.e. neurons having more than one applicable rule, since we only consider deterministic systems in this work. Minimization of the number of neurons of a $\overline{\Pi}$ simulating a Π is also desirable, including providing bounds to the number of neurons, spikes, and types of spiking rules. We will also use the matrix representation of SNP systems without delays from [15] and then use massively parallel processors such as graphics processing units to create simulations of computations as was done in [2]. Lastly, certain results and applications of SNP systems that use delays can be converted to SNP systems without delays e.g. generating automatic sequences as in [4], and performing arithmetic operations as in [16], among others.

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